



University of
Strathclyde
Science

Lecture 11: Languages and Recursion

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CS103 Machines, Languages and Computation
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The First Half of the Class

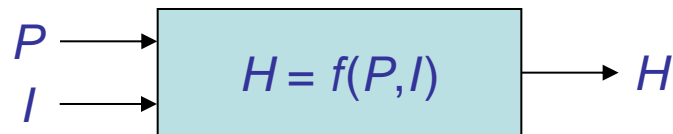
- The first half of the class was about the fact that all computers are equivalent and that non-computable functions exist:

If P is the set of all computer programs and F is the set of all functions $f: n \rightarrow m$ such that n, m are members of the set of natural numbers N , then $|F| > |P|$ and therefore there are some functions in F for which no computer program can exist.

- We now know why this is: P is infinite but countable, and F is infinite and uncountable.

The Halting Problem

- One of the most famous non-computable functions is called “The Halting Problem”
- Given a computer program, P , and a set of inputs to that program, I , decide (in finite time) whether or not the program halts or runs forever:



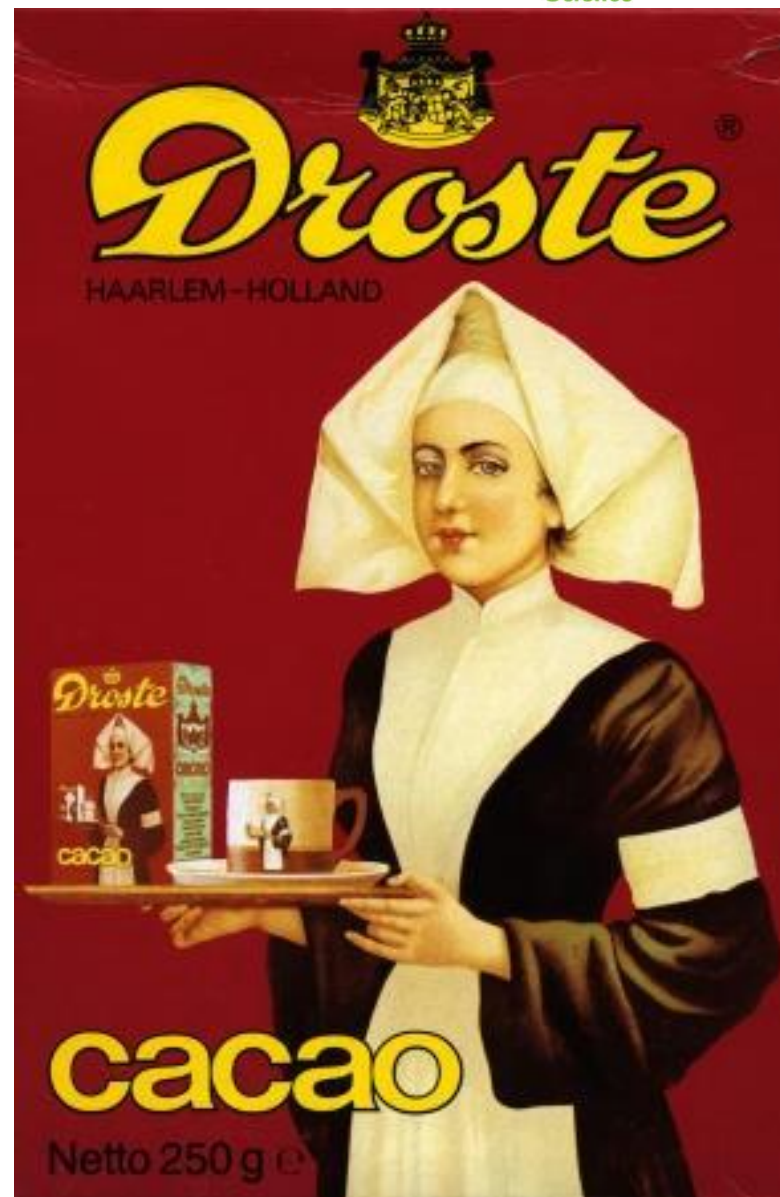
- $H = 1$ means the program halts, $H = 0$ means that it runs forever
- It can be shown that no program can exist for $f(P, I)$

The Second Half of the Class

- Week 7: Recursion and Languages: how to use recursion to allow for potentially unbounded sentences.
- Week 8: Recursion and Strings: how to specify an infinite set of strings using a finite set of rules.
- Week 9: Recursion and Logical Definitions: what does it mean to be someone's ancestor?
- Week 10: The λ -calculus: how can we create and use computable functions using only symbols?
- Week 11: Recursion and Function Definitions: how can we use recursion to create functions?

What is Recursion?

- Recursion is a method of defining structures in which the structure being defined may be used within its own definition.
- The term is also used more generally to describe a process of repeating objects in a self-similar way.
- Recursion leads to “nested” structures. Sometimes the nature of the nesting is clear, but often we have to look hard to find it and encode the recursive process.
- Recursion can be found in many places: in languages, logical definitions, data structures, programs, ...



Recursion by Example

- How do we define the members of the set of natural numbers, $N = \{0, 1, 2, 3, \dots\}$?
- We can do it by recursion:
 - 0 is a member of N
 - If n is a member of N , $n+1$ is a member of N
- We need *both* statements: either one on its own is not enough to complete the definition.
- Notice the very strong similarity between this definition and proof by induction, as introduced in Lecture 4.

What are Languages?

- Languages are devices used to transfer information between people, and between people and machines
- Statements in a language consist of sequences of symbols, such as “John ate a frog” or “let $x = \sqrt{k} + 1$ ”
- Languages have *structure*: we can’t just produce the symbols in a random order
- We can *recognise* legal and illegal sentences:

John ate a big frog. 😊

big ate John frog a. 😞

let $x = \sqrt{+} * \text{if}(k)$ 😞

Syntax and Semantics

- Languages have *structure*: we can't just produce the symbols in a random order
- The structure of a language is called *syntax* and can be specified using a *grammar*
- The meaning of an expression of language is called the *semantics* of the expression (e.g. “let x=3” is an instruction to set the value of variable x to be 3)
- For the moment, let's concentrate on the structure of language and look at some of the repeating structures we can find...

Types of Languages

- Natural languages which have evolved for human communication: English, French, Italian, Chinese.
- Invented languages for communication: Esperanto, Loglan, Klingon.
- Computer languages: Java, Lisp, C, Algol, Fortran, Basic, Cobol, Pascal, ML, Haskell, Perl, Python, ...
- Formal languages: mathematical logic
- Invented languages, e.g. $L = a^n b^n c^n$ – used to find out what languages we can describe with our grammars

Context-Free Grammars

- Most computer languages can be specified using context-free grammars
- Rules are of the form:

Symbol \rightarrow list of Symbols and Terminal Symbols

e.g.

Exp \rightarrow “if” Cond “then” Exp

Exp \rightarrow “print” Var

Cond \rightarrow Var “=” Value

Var \rightarrow “a” | “b” | “c” | ... | “y” | “z” (| means “or”)

Value \rightarrow “0” | “1”

Coping with Repeating Structure

Exp \rightarrow “if” Cond “then” Exp

Exp \rightarrow “print” Var

Cond \rightarrow Var “=” Value

Var \rightarrow “x” | “y” | “z”

Value \rightarrow “0” | “1”

- This grammar can generate repeating structures:

if a=0 then

 if b=1 then

 if c=0 then

 if d=0 then ...

Recursion

- Recursion occurs when you define something in terms of itself, e.g. the rule “Exp \rightarrow “if” Cond “then” Exp”
- From GEB: “Hofstadter’s Law: it always takes longer than you expect, even when you take into account Hofstadter’s law.”
- Python factorial function:

```
def factorial (n):  
    if n == 0:  
        return 1  
    else:  
        return n * factorial (n - 1)
```

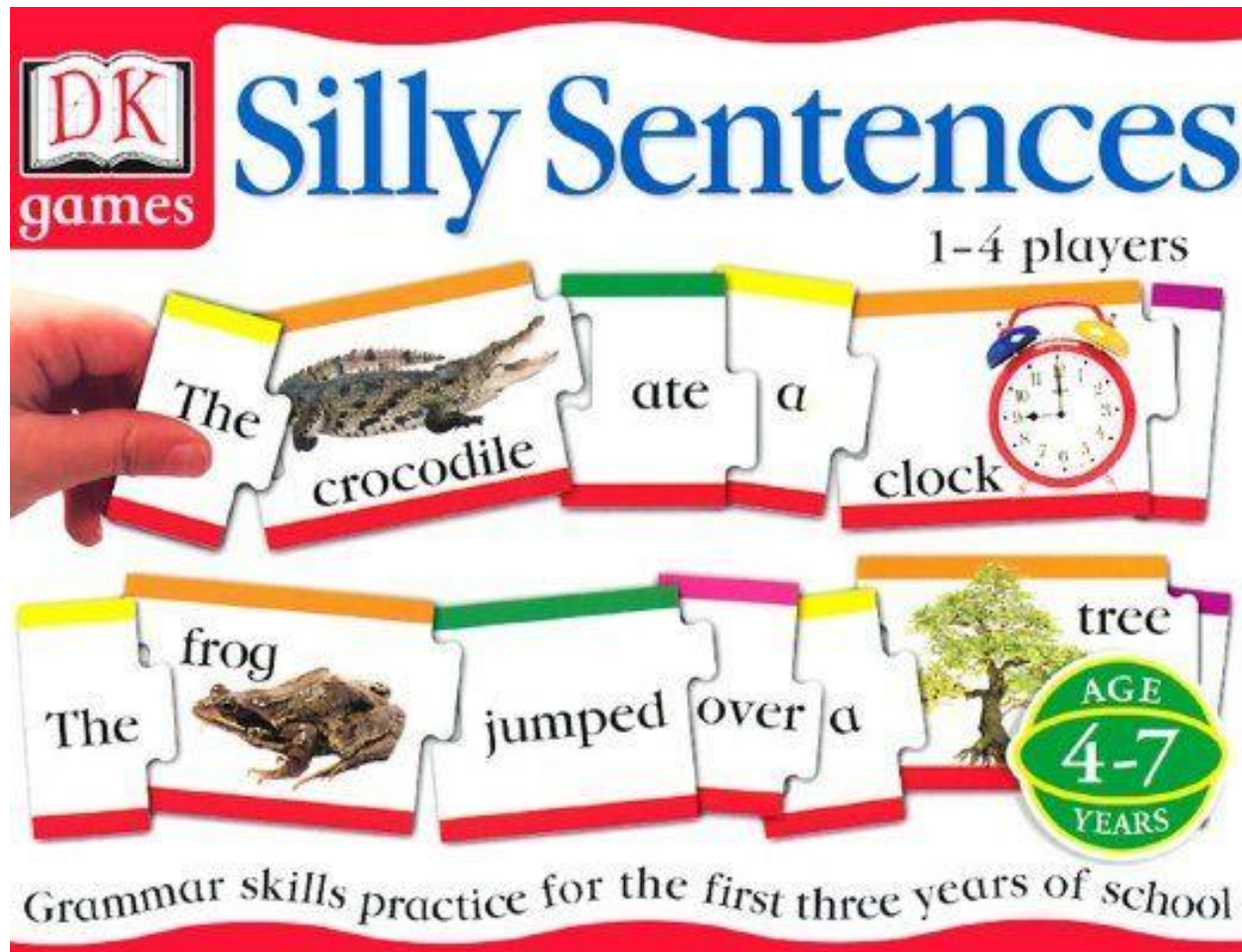
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Silly Sentences



Silly Sentences

- This is a game by Dorling Kindersley containing jigsaw pieces with words on them.
- On the box, it says “Grammar skills practice for the first three years at school”.
- You can assemble the pieces to create random (but grammatical) sentences, such as:

The yellow crocodile ate a rubber clock.

The old fat hairy princess snores.

Brian ate a big purple frog.

Silly Sentences

- Can we write a program to generate legal sentences?
- We can make use of production rules, similar to the rules of the MIU system:

Sentence \rightarrow NounPhrase Verb1

NounPhrase \rightarrow Name

Name \rightarrow “Brian”

Name \rightarrow “Beryl”

Verb1 \rightarrow “snores”

Verb1 \rightarrow “smells”

- Words in “...” are called *terminal symbols*

Sentence Generation Algorithm

- Start with “Sentence”.
- Find the left-most non-terminal symbol.
- Expand it using a random rule that applies.
- Repeat from 2 until all the symbols are terminals.

“Sentence” → “NounPhrase Verb1”

“NounPhrase Verb1” → “Name Verb1”

“Name Verb1” → “Brian Verb1”

“Brian Verb1” → “Brian snores”

A Larger Grammar

Sentence \rightarrow NounPhrase Verb2 NounPhrase

NounPhrase \rightarrow Name

NounPhrase \rightarrow Article NounGroup

NounGroup \rightarrow Noun

NounGroup \rightarrow Adjective Noun

Name \rightarrow “Brian” | “Beryl”

Verb2 \rightarrow “ate” | “fed” | “chased” | “kissed”

Article \rightarrow “a” | “the”

Adjective \rightarrow “big” | “purple” | “ugly” | “fat” | “eccentric”

Noun \rightarrow “frog” | “dog” | “clock” | “monster” | “princess”

Derivation of a Sentence

- We can derive sentences with grammars just like we did with the MIU-system:

Sentence

NounPhrase Verb2 NounPhrase

Name Verb2 NounPhrase

“Brian” Verb2 NounPhrase

“Brian” “ate” NounPhrase

“Brian” “ate” Article NounGroup

“Brian” “ate” “a” NounGroup

“Brian” “ate” “a” Adjective Noun

“Brian” “ate” “a” “big” Noun

“Brian” “ate” “a” “big” “frog” ■

Assignment 6

- Generate some random sentences using the grammar.
- The grammar can generate “Brian ate a big frog” or “Brian ate a purple frog”.
- But I want it to generate “Brian ate a big purple frog”.
- Or even “Brian ate a big fat ugly purple eccentric frog”.
- Can you modify my grammar so it can generate noun phrases containing any number of adjectives, without increasing the number of rules?
- Reading: Chapter 5 of Gödel, Escher, Bach.