

The Sub-Additives:

A Proof Theory for Probabilistic Choice extending Linear Logic

Fundamental Structure of Computation and Deduction (FSCD'19)

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Lets talk about Coffee



Processes as Formulae

In process calculi:

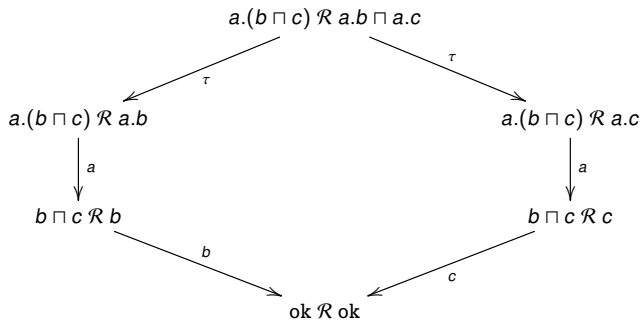
$a.(b \sqcap c)$ simulates $a.b \sqcap a.c$

In MAV:

$\vdash a \triangleleft (b \& c) \multimap (a \triangleleft b) \& (a \triangleleft c)$

Processes as Formulae

In process calculi:



In MAV:

$$\frac{\frac{\frac{\vdash \bar{a}, a}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), a \triangleleft b} \quad \frac{\frac{\vdash \bar{b}, b}{\vdash \bar{b} \oplus \bar{c}, b}}{\vdash \bar{a}, a \quad \vdash \bar{b} \oplus \bar{c}, b}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), a \triangleleft b} \quad \frac{\frac{\frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{a}, a \quad \vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), a \triangleleft c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), (a \triangleleft b) \& (a \triangleleft c)}$$

From Proofs in MAV to Simulations

$$\frac{\frac{\frac{\vdash \bar{a}, a}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), a \triangleleft b} \quad \frac{\frac{\vdash \bar{b}, b}{\vdash \bar{b} \oplus \bar{c}, b}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), a \triangleleft b}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), (a \triangleleft b) \& (a \triangleleft c)} \quad \frac{\frac{\frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), a \triangleleft c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), (a \triangleleft b) \& (a \triangleleft c)}}$$

From Proofs in MAV to Simulations

$$\frac{\frac{\frac{\vdash \bar{a}, a}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft b}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, (a \triangleleft b) \& (a \triangleleft c)}{\frac{\frac{\frac{\vdash \bar{b}, b}{\vdash \bar{b} \oplus \bar{c}, b}}{\vdash \bar{a}, a} \quad \frac{\frac{\frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{a}, a} \quad \frac{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, (a \triangleleft b) \& (a \triangleleft c)}}$$

By cut elimination.

$$\frac{\frac{\frac{\vdash \bar{a}, a}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft b}{\vdash \bar{a}, a} \quad \frac{\frac{\frac{\vdash \bar{b}, b}{\vdash \bar{b} \oplus \bar{c}, b}}{\vdash \bar{a}, a} \quad \frac{\frac{\frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{a}, a} \quad \frac{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft b}}{\vdash \bar{a}, a} \quad \frac{\frac{\frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{a}, a} \quad \frac{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft c}}$$

From Proofs in MAV to Simulations

$$\frac{\frac{\frac{\vdash \bar{a}, a}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft b}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, (a \triangleleft b) \ \& \ (a \triangleleft c)}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, (a \triangleleft b) \ \& \ (a \triangleleft c)}{\frac{\frac{\frac{\vdash \bar{b}, b}{\vdash \bar{b} \oplus \bar{c}, b}}{\vdash \bar{a}, a} \quad \frac{\frac{\frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{a}, a} \quad \frac{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, (a \triangleleft b) \ \& \ (a \triangleleft c)}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, (a \triangleleft b) \ \& \ (a \triangleleft c)}$$

By cut elimination.

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Removing left-most interacting atoms.

$$\frac{\frac{\vdash \bar{b}, b}{\vdash \bar{b} \oplus \bar{c}, b}}{\vdash \bar{b} \oplus \bar{c}, b} \quad \frac{\frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{b} \oplus \bar{c}, c}$$

From Proofs in MAV to Simulations

$$\frac{\frac{\frac{\vdash \bar{a}, a}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft b}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, (a \triangleleft b) \& (a \triangleleft c)}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, (a \triangleleft b) \& (a \triangleleft c)}{\frac{\frac{\frac{\vdash \bar{b}, b}{\vdash \bar{b} \oplus \bar{c}, b}}{\vdash \bar{a}, a \quad \frac{\vdash \bar{b}, b}{\vdash \bar{b} \oplus \bar{c}, b}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft b} \quad \frac{\frac{\frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{a}, a \quad \frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, (a \triangleleft b) \& (a \triangleleft c)}}$$

By cut elimination.

$$\frac{\frac{\frac{\vdash \bar{a}, a}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft b}{\vdash \bar{a}, a \quad \frac{\vdash \bar{b}, b}{\vdash \bar{b} \oplus \bar{c}, b}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft b} \quad \frac{\frac{\frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{a}, a \quad \frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft c}}{\vdash \bar{a}, a \quad \frac{\vdash \bar{b}, b}{\vdash \bar{b} \oplus \bar{c}, b}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft b} \quad \frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c})}, a \triangleleft c}$$

Removing left-most interacting atoms.

$$\frac{\frac{\vdash \bar{b}, b}{\vdash \bar{b} \oplus \bar{c}, b}}{\vdash \bar{b}, b} \quad \frac{\frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{c}, c}}$$

Removing left-most interacting atoms.

$\vdash \circ, \circ$

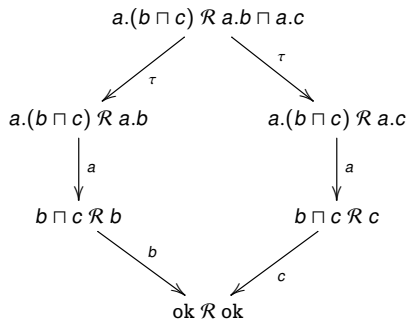
Processes as Formulae

$$\frac{\frac{\frac{\vdash \bar{b}, b}{\vdash \bar{b} \oplus \bar{c}, b}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), a \triangleleft b}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), (a \triangleleft b) \& (a \triangleleft c)}}{\frac{\frac{\frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), a \triangleleft c}}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), (a \triangleleft b) \& (a \triangleleft c)}}$$

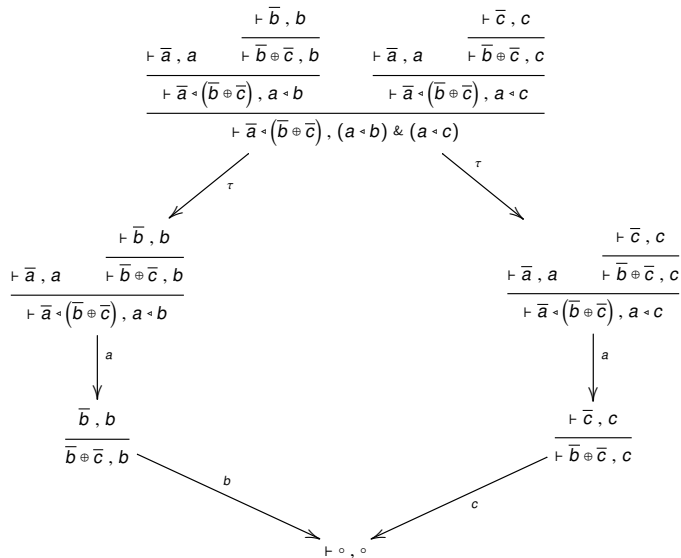
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$$\frac{\bar{b}, b}{\bar{b} \oplus \bar{c}, b} \qquad \frac{\vdash \bar{c}, c}{\vdash \bar{b} \oplus \bar{c}, c}$$

$\vdash \circ, \circ$



Processes as Formulae



Theorem (soundness)

If $\vdash \llbracket P \rrbracket \multimap \llbracket Q \rrbracket$ then P simulates Q .

Probabilistic Tests (already probabilistic)

$a.b \sqcap a.c$ does not simulate $a.(b \sqcap c)$

Consider probabilistic test:

$$\bar{a}.\left(\bar{b}.\omega +_{1/2} \bar{c}.\omega\right)$$

Probabilistic Tests (already probabilistic)

$a.b \sqcap a.c$ does not simulate $a.(b \sqcap c)$

Consider probabilistic test:

$$\bar{a} . (\bar{b} . \omega +_{1/2} \bar{c} . \omega)$$

passes 50%

v.s.

passes 100%

Consequently:

$$\not\vdash (a \triangleleft b) \& (a \triangleleft c) \multimap a \triangleleft (b \& c)$$

In process calculi:

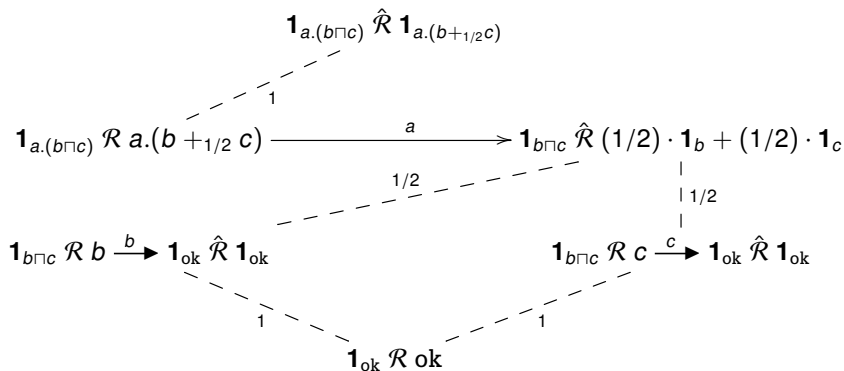
$$a.(b \sqcap c) \quad \textit{simulates} \quad a.(b +_{1/2} c)$$

In Δ MAV:

$$\vdash a \triangleleft (b \& c) \multimap a \triangleleft (b \oplus_{1/2} c)$$

Probabilistic Processes

In process calculi:



In ΔMAV :

$$\vdash a \triangleleft (b \& c) \dashv\vdash a \triangleleft (b \oplus_{1/2} c)$$

In process calculi:

$$a.(b \sqcap c) \quad \textit{simulates} \quad a.(b +_{1/2} c)$$

In Δ MAV:

$$\frac{\vdash \bar{a}, a \quad \vdash \bar{b} \oplus \bar{c}, b \oplus_{1/2} c}{\vdash \bar{a} \triangleleft (\bar{b} \oplus \bar{c}), a \triangleleft (b \oplus_{1/2} c)}$$

Probabilistic Processes

In process calculi:

$a.(b \sqcap c)$ *simulates* $a.(b +_{1/2} c)$

In Δ MAV:

$$\frac{\frac{\frac{\vdots}{(\bar{b} \oplus \bar{c}) \wp (b \oplus_{1/2} c)}}{(\bar{a} \wp a) \triangleleft ((\bar{b} \oplus \bar{c}) \wp (b \oplus_{1/2} c))}}{(\bar{a} \triangleleft (\bar{b} \oplus \bar{c})) \wp (a \triangleleft (b \oplus_{1/2} c))}$$

$$\circ \triangleleft P \equiv P$$

$$\frac{C\{\circ\}}{C\{\bar{a} \wp a\}}$$

$$\frac{C\{(P \wp R) \triangleleft (Q \wp S)\}}{C\{(P \triangleleft Q) \wp (R \triangleleft S)\}}$$

Probabilistic Processes

In process calculi:

$a.(b \sqcap c)$ *simulates* $a.(b +_{1/2} c)$

In Δ MAV:

$$\frac{\frac{\vdots}{\frac{((\bar{b} \oplus \bar{c}) \&_{1/2} (\bar{b} \oplus \bar{c})) \wp (b \oplus_{1/2} c)}{(\bar{b} \oplus \bar{c}) \wp (b \oplus_{1/2} c)}}{(\bar{a} \wp a) \triangleleft ((\bar{b} \oplus \bar{c}) \wp (b \oplus_{1/2} c))}}{(\bar{a} \triangleleft (\bar{b} \oplus \bar{c})) \wp (a \triangleleft (b \oplus_{1/2} c))}$$

$$P \equiv P \&_p P$$

Probabilistic Processes

In process calculi:

$a.(b \sqcap c)$ *simulates* $a.(b +_{1/2} c)$

In Δ MAV:

$$\begin{array}{c}
 \vdots \\
 \frac{\frac{((\bar{b} \oplus \bar{c}) \wp b) \&_{1/2} ((\bar{b} \oplus \bar{c}) \wp c)}{((\bar{b} \oplus \bar{c}) \&_{1/2} (\bar{b} \oplus \bar{c})) \wp (b \oplus_{1/2} c)}}{(\bar{b} \oplus \bar{c}) \wp (b \oplus_{1/2} c)}}{\frac{(\bar{a} \wp a) \triangleleft ((\bar{b} \oplus \bar{c}) \wp (b \oplus_{1/2} c))}{(\bar{a} \triangleleft (\bar{b} \oplus \bar{c})) \wp (a \triangleleft (b \oplus_{1/2} c))}}
 \end{array}
 \qquad
 \frac{C\{(P \wp R) \&_p (Q \wp S)\}}{C\{(P \&_p Q) \wp (R \oplus_p S)\}}$$

$$P \equiv P \&_p P$$

Probabilistic Processes

In process calculi:

$$a.(b \sqcap c) \quad \text{simulates} \quad a.(b +_{1/2} c)$$

In Δ MAV:

$$\frac{\frac{\frac{\frac{\frac{\frac{\frac{\circ}{\circ \&_{1/2} \circ}}{(\bar{b} \&_{1/2} b) \&_{1/2} (\bar{c} \&_{1/2} c)}}{((\bar{b} \oplus \bar{c}) \&_{1/2} b) \&_{1/2} ((\bar{b} \oplus \bar{c}) \&_{1/2} c)}}{((\bar{b} \oplus \bar{c}) \&_{1/2} (\bar{b} \oplus \bar{c})) \&_{1/2} (b \oplus_{1/2} c)}}{(\bar{b} \oplus \bar{c}) \&_{1/2} (b \oplus_{1/2} c)}}{(\bar{a} \&_{1/2} a) \triangleleft ((\bar{b} \oplus \bar{c}) \&_{1/2} (b \oplus_{1/2} c))}}{(\bar{a} \triangleleft (\bar{b} \oplus \bar{c})) \&_{1/2} (a \triangleleft (b \oplus_{1/2} c))}}{\frac{C\{\circ\}}{C\{\bar{a} \&_{1/2} a\}}}$$

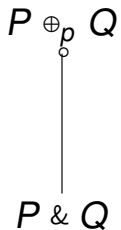
$$\frac{C\{P\}}{C\{P \oplus Q\}}$$

$$\frac{C\{Q\}}{C\{P \oplus Q\}}$$

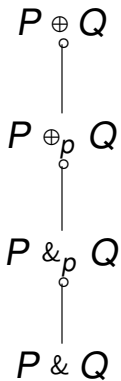
$$\frac{C\{(P \&_p R) \&_p (Q \&_{1-p} S)\}}{C\{(P \&_p Q) \&_{1-p} (R \oplus_p S)\}}$$

$$P \equiv P \&_p P$$

The Sub-Additives



The Sub-Additives



Distributivity Properties

In process calculi:

$$(a \parallel c) +_{1/2} (b \parallel d) \quad \text{is unrelated to} \quad (a +_{1/2} b) \parallel (c +_{1/2} d)$$

In Δ MAV:

$$(a \otimes c) \oplus_{1/2} (b \otimes d) \quad \text{is unrelated to} \quad (a \oplus_{1/2} b) \otimes (c \oplus_{1/2} d)$$

Distributivity Properties

In process calculi:

$$(a \sqcap b) \oplus_{1/2} (a \sqcap c) \quad \text{simulates} \quad a \sqcap (b \oplus_{1/2} c)$$

In Δ MAV:

$$\vdash (a \& b) \oplus_{1/2} (a \& c) \multimap a \& (b \oplus_{1/2} c)$$

Proof:

$$\begin{array}{c}
 \frac{\circ}{\circ \& \circ} \text{ tidy} \\
 \frac{\circ \& \circ}{\circ \& \circ} \text{ idempotency} \\
 \frac{\circ \& \circ}{\circ \& \circ} \text{ by atomic interaction} \\
 \frac{\circ \& \circ}{\circ \& \circ} \text{ by confine} \\
 \frac{\circ \& \circ}{\circ \& \circ} \text{ idempotency} \\
 \frac{\circ \& \circ}{\circ \& \circ} \text{ by choose} \\
 \frac{\circ \& \circ}{\circ \& \circ} \text{ by external}
 \end{array}$$

Conclusions

Sub-additives arise from the probabilistic content of linear logic.

Further remarks (in paper):

- ▶ Preserved in all contexts.
- ▶ Permits action refinement.

$$\not\vdash a \triangleleft a \multimap a \otimes a \quad (\text{no auto-concurrency})$$

- ▶ **Medial rules** necessary for cut elimination.

$$\frac{(P \sqcap R) \sqcup (Q \sqcap S)}{(P \sqcup Q) \sqcap (R \sqcup S)} \text{medial} \quad \text{where } (\sqcap, \sqcup) \in \{ (\wp, \oplus_q), (\&_p, \oplus_q), (\&, \oplus_q), (\&, \&_p), (\wp, \triangleleft), (\&, \triangleleft), (\&_p, \triangleleft), (\triangleleft, \oplus_q) \}$$

- ▶ Cut elimination demands **decomposition**, based on the topology of proofs.

Questions (for audience):

- ▶ Do sub-additives arise in other semantics?