THE SIGMA-SEQUENCE AND COUNTING OCCURRENCES OF SOME PATTERNS, SUBSEQUENCES AND SUBWORDS

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ABSTRACT. We consider *sigma-words*, which are words used by Evdokimov in the construction of the sigma-sequence [7]. We then find the number of occurrences of certain patterns, subsequences and subwords in these words.

1. INTRODUCTION AND BACKGROUND

We write permutations as words $\pi = a_1a_2\cdots a_n$, whose letters are distinct and usually consist of the integers $1, 2, \ldots, n$.

An occurrence of a pattern p in a permutation π is "classically" defined as a subsequence in π (of the same length as the length of p) whose letters are in the same relative order as those in p. Formally speaking, for $r \leq n$, we say that a permutation σ in the symmetric group $\mathcal{S}_n$ has an occurrence of the pattern $p \in \mathcal{S}_r$ if there exist $1 \leq i_1 < i_2 < \cdots < i_r \leq n$ such that $p = \sigma(i_1)\sigma(i_2)\ldots\sigma(i_r)$ in reduced form. The *reduced form* of a permutation σ on a set $\{j_1, j_2, \ldots, j_r\}$, where $j_1 < j_2 < \cdots < j_r$, is a permutation σ_1 obtained by renaming the letters of the permutation σ so that j_i is renamed i for all $i \in \{1, \ldots, r\}$. For example, the reduced form of the permutation 3651 is 2431. The first case of classical patterns studied was that of permutations avoiding a pattern of length 3 in $\mathcal{S}_3$. Knuth [13] found that, for any $\tau \in \mathcal{S}_3$, the number $|\mathcal{S}_n(\tau)|$ of n-permutations avoiding τ is C_n, the nth Catalan number. Later, Simion and Schmidt [16] determined the number $|\mathcal{S}_n(P)|$ of permutations in $\mathcal{S}_n$ simultaneously avoiding any given set of patterns $P \subseteq \mathcal{S}_3$.

In [1] Babson and Steingrímsson introduced *generalized permutation patterns* that allow the requirement that two adjacent letters in a pattern must be adjacent in the permutation. In order to avoid confusion we write a "classical" pattern, say 231, as 2-3-1, and if we write, say 2-31, then we mean that if this pattern occurs in the permutation, then the letters in the permutation that correspond to 3 and 1 are adjacent. For example, the permutation $\pi = 516423$ has only one occurrence of the pattern 2-31, namely the subword 564, whereas the pattern 2-3-1 occurs, in addition, in the subwords 562 and 563. A motivation for introducing these patterns in [1] was the study of Mahonian statistics. A number of interesting results on generalized patterns were obtained in [6]. Relations to several well studied combinatorial structures, such as set partitions, Dyck paths, Motzkin paths and involutions, were shown there.

Burstein [2] considered words instead of permutations. In particular, he found the number $|[k]^n(P)|$ of words of length n in a k-letter alphabet that avoid all patterns from a set $P \subseteq \mathcal{S}_3$ simultaneously. Burstein and Mansour [3] (resp. [4, 5]) considered forbidden patterns (resp. generalized patterns) with repeated letters.

The most attention, in the papers on classical or generalized patterns, is paid to finding exact formulas and/or generating functions for the number of words or permutations avoiding, or having k occurrences of, certain patterns. In [11] the authors suggested another problem, namely counting the number of occurrences of certain patterns in certain words. These words were chosen to be the set of all finite approximations of a sequence generated by a *morphism* with certain restrictions. A motivation for this choice was the interest in studying classes of sequences and words that are defined by iterative schemes [14, 15]. In [12] the authors also studied the number of occurrences of certain patterns in certain words. But there they choose these words to be the subdivision stages from which the *Peano curve* is obtained. The authors called these words the *Peano words.* The Peano curve was studied by the Italian mathematician Giuseppe Peano in 1890 as an example of a continuous space filling curve.

In the present paper we consider the *sigma-words*, which are words used by Evdokimov in construction of the *sigma-sequence* [7]. Evdokimov used this sequence to construct chains of maximal length in the n-dimensional unit cube. Independent interest to the sigma-sequence appears in connection with the well-known *Dragon curve*, discovered by physicist John E. Heighway and defined as follows: we fold a sheet of paper in half, then fold in half again, and again, etc. and then unfold in such way that each crease created by the folding process is opened out into a 90-degree angle. The "curve" refers to the shape of the partially unfolded paper as seen edge on. If one travels along the curve, some of the creases will represent turns to the left and others turns to the right. Now if 1 indicates a turn to the right, and 2 to the left, and we start travelling along the curve indicating the turns, we get the sigma-sequence [8]. In [10] the sigma-sequence was studied from another point of view. It was proved there that this sequence cannot be defined by iterated morphism.

Since the sigma-sequence w_σ is a sequence in a 2-letter alphabet, we consider patterns in 2-letter alphabets. Moreover, the patterns in a 1-letter alphabet (for example 1-1-1) correspond to two subsequences (for this example, these subsequences are 1-1-1 and 2-2-2), whereas the patterns in a 2-letter alphabet (with at least one letter 2) uniquely determine the subsequences in w_σ that correspond to them, and conversely. For example, an occurrence of the pattern 1-2-1 is an occurrence of the subsequence 1-2-1, whereas an occurrence of the subsequence (subword) 211 is an occurrence of the pattern 211. Thus, any of our results for a pattern, can be interpreted in term of subsequences or subwords, depending on the context, and conversely.

In our paper we give either an explicit formula or recurrence relation for the number of occurrences for some classes of patterns, subwords and subsequences in the sigma-words. In particular, Theorem 4, allows to find the number of occurrences of an arbitrary generalized pattern without internal dashes of length ℓ, provided we know four certain numbers that can be easily calculated for the sigma-words C_k, D_k, C_{k+1} and D_{k+1} (to be defined below), where $k = \lceil \log_2 \ell \rceil$. Theorem 9 gives a recurrence relation for counting occurrences of patterns of the form τ_1-τ_2. In Section 6 we discuss occurrences of patterns of the form τ_1-τ_2-$\cdots$-τ_k, where the pattern τ_i does not overlap with the patterns τ_{i-1} and τ_{i+1} for $i = 1, 2, \ldots, k-1$. Finally, Section 7 deals with patterns of the form $[\tau_1$-τ_2-$\cdots$-$\tau_k]$, $[\tau_1$-τ_2-$\cdots$-$\tau_k)$ and $(\tau_1$-τ_2-$\cdots$-$\tau_k]$ in Babson and Steingrímsson notation, which means that we use "$[x$" in a pattern p to indicate that in an occurrence of p, the letter corresponding to the x must be the first letter of a word under consideration, whereas if we use "$y]$", we mean that the letter corresponding to y must be the last (rightmost) letter in the word.

2. Preliminaries

In [7], Evdokimov constructed chains of maximal length in the n-dimensional unit cube using the *sigma-sequence.* The sigma-sequence w_σ was defined there by the following inductive scheme:

$$C_1 = 1, \qquad D_1 = 2$$
$$C_{k+1} = C_k 1 D_k, \qquad D_{k+1} = C_k 2 D_k$$
$$k = 1, 2, \ldots$$

and $w_\sigma = \lim_{k\to\infty} C_k$. Thus, the initial letters of w_σ are 11211221112212.... We call the words C_k the *sigma words.* The first four values of the sequence $\{C_k\}_{k\geq 1}$ are 1, 112, 1121122, 112112211122122.

In [10] an equivalent definition of w_σ was given: any natural number $n \neq 0$ can be presented unambiguously as $n = 2^t(4s + \sigma)$, where $\sigma < 4$, and t is the greatest natural number such that 2^t divides n. If n runs through the natural numbers then σ runs through some sequence consisting of 1 and 3. If we substitute 3 by 2 in this sequence, we get w_σ.

In this paper we count occurrences of patterns in the sigma-words, which are particular initial subwords of w_σ. However, the challenging question is to find the number of occurrences of patterns or subwords in an arbitrary initial subword of w_σ, or more generally, in a subword of w_σ starting in the position i and ending in the position j.

It turns out that for counting occurrences of certain patterns or subwords in C_n, one needs to know the number of occurrences of certain patterns in D_n. So, in the most cases, we give results for both C_n and D_n. However, our main purpose is the words C_n for $n \geq 1$, and in some propositions and examples we do not consider D_n.

In what follows, we give initial values for the words C_i and D_i:

$$\begin{array}{ll}
C_1 = 1 & D_1 = 2 \\
C_2 = 112 & D_2 = 122 \\
C_3 = 1121122 & D_3 = 1122122 \\
C_4 = 112112211122122 & D_4 = 112112221122122 \\
C_5 = 1121122111221221112112221122122 & D_5 = 1121122111221222112112221122122
\end{array}$$

We now give some other definitions.

A *descent* (resp. *rise*) in a word $w = a_1a_2\dots a_n$ is an i such that $a_i > a_{i+1}$ (resp. $a_i < a_{i+1}$). It follows from the definitions that an occurrence of a descent (resp. rise) is an occurrence of the pattern 21 (resp. 12).

Let c_n^τ (resp. d_n^τ) denote the number of occurrences of the pattern τ in C_n (resp. D_n).

Suppose a word $W = AaB$, where A and B are some words of the same length, and a is a letter. We define the *kernel of order k* for the word W to be the subword consisting of the $k-1$ rightmost letters of A, the letter a, and the $k-1$ leftmost letters of B. We denote it by $\mathcal{K}_k(W)$. For example, $\mathcal{K}_3(111211221) = 12112$. If $|A| < k-1$ then we assume $\mathcal{K}_k(W) = \epsilon$, that is, the kernel in this case is the empty word. Also, $m_k(\tau, W)$ denotes the number of occurrences of the pattern (or the word, or the subsequence depending on the context) τ in $\mathcal{K}_k(W)$.

We denote $x\text{-}x\text{-}\cdots\text{-}x$ (ℓ times) by x^ℓ. Also, $\lceil a \rceil$ denotes the least natural number b such that $a \le b$.

3. PATTERNS 1-1-$\cdots$-1, 1-2 AND 2-1

It is easy to see that $|C_n| = |D_n| = 2^n - 1$. The following lemma gives the number of the letters 1 and 2 in C_n and D_n.

Lemma 1. *The number of 1s (resp. 2s) in C_n is 2^{n-1} (resp. $2^{n-1}-1$). The number of 1s (resp. 2s) in D_n is $2^{n-1}-1$ (resp. 2^{n-1}).*

Proof. It is enough to find the number of 1s c_n and d_n in C_n and D_n respectively, since the number of 2s in C_n and D_n are obviously equal to $|C_n| - c_n$ and $|D_n| - d_n$ respectively.

It is easy to see from the structure of C_n and D_n that

$$\begin{cases} c_n = c_{n-1} + d_{n-1} + 1, \\ d_n = c_{n-1} + d_{n-1}, \end{cases}$$

together with $c_1 = 1$ and $d_1 = 0$. The solution to this recurrence is $c_n = 2^{n-1}$ and $d_n = 2^{n-1}-1$. ☐

Proposition 2. *The number occurrences of the subsequence 1^k (resp. 2^k) in C_n is $\binom{2^{n-1}}{k}$ (resp. $\binom{2^{n-1}-1}{k}$). Thus, the number of occurrences of the pattern 1^k in C_n is equal to*

$$c_n^{1^k} = \binom{2^{n-1}}{k} + \binom{2^{n-1}-1}{k} = \frac{2^n - k}{2^{n-1}-k}\binom{2^{n-1}-1}{k}.$$

Proof. From Lemma 1, there are 2^{n-1} (resp. $2^{n-1}-1$) occurrences of the letter 1 (resp. 2) in C_n, and thus there are $\binom{2^{n-1}}{k}$ (resp. $\binom{2^{n-1}-1}{k}$) occurrences of the subsequence 1^k (resp. 2^k) there. ☐

Proposition 3. *We have that for all $n \ge 2$, $c_n^{1\text{-}2} = d_n^{1\text{-}2} = 2 \cdot 4^{n-2} + (n-2) \cdot 2^{n-2}$, and $c_n^{2\text{-}1} = d_n^{2\text{-}1} = 2 \cdot 4^{n-2} - n \cdot 2^{n-2}$.*

Proof. Let us first consider the pattern 1-2. An occurrence of this pattern in $C_n = C_{n-1}1D_{n-1}$ is either inside C_{n-1}, or inside D_{n-1}, or the letter 1 is from the word $C_{n-1}1$, whereas the letter 2 is from the word D_{n-1}. Thus

$$c_n^{1\text{-}2} = c_{n-1}^{1\text{-}2} + d_{n-1}^{1\text{-}2} + \{\text{ (the number of 1s in } C_{n-1}) + 1\} \cdot \{\text{ the number of 2s in } D_{n-1}\}.$$

Using the same considerations for $D_n = C_{n-1}2D_{n-1}$, one can get

$$d_n^{1\text{-}2} = c_{n-1}^{1\text{-}2} + d_{n-1}^{1\text{-}2} + \{\text{ the number of 1s in } C_{n-1}\} \cdot \{\text{ (the number of 2s in } D_{n-1}) + 1\}.$$

The number of 1s and 2s in C_{n-1} and D_{n-1} is given in Lemma 1. So,

$$\begin{cases} c_n^{1\text{-}2} = c_{n-1}^{1\text{-}2} + d_{n-1}^{1\text{-}2} + 2^{n-2} \cdot (2^{n-2}+1) \\ d_n^{1\text{-}2} = c_{n-1}^{1\text{-}2} + d_{n-1}^{1\text{-}2} + 2^{n-2} \cdot (2^{n-2}+1) \end{cases} \Leftrightarrow$$

$$\begin{pmatrix} c_n^{1\text{-}2} \\ d_n^{1\text{-}2} \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} c_{n-1}^{1\text{-}2} \\ d_{n-1}^{1\text{-}2} \end{pmatrix} + \begin{pmatrix} 2^{n-2} \cdot (2^{n-2}+1) \\ 2^{n-2} \cdot (2^{n-2}+1) \end{pmatrix} \tag{1}$$

together with $c_2^{1\text{-}2} = 2$ and $d_2^{1\text{-}2} = 2$. Here, and several times in what follows, we need to solve recurrence relations of the form

$$x_n = Ax_{n-1} + b,$$

where A is a matrix, and x_n, x_{n-1} and b are some vectors, where b sometimes depends on n. We recall from linear algebra that such relations can be solved by diagonalization of the matrix A, that is, by writing $A = VDV^{-1}$, where D is a diagonal matrix consisting of eigenvalues of A, and the columns of V are eigenvectors of A. For example, if A is a 2×2 matrix that consists of 1s, then we use

$$\begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ -1 & 1 \end{pmatrix} \begin{pmatrix} 0 & 0 \\ 0 & 2 \end{pmatrix} \begin{pmatrix} 1/2 & -1/2 \\ 1/2 & 1/2 \end{pmatrix}$$

for computing powers of A, and thus for solving the recurrence relations. For the recurrence 1, we get that for all $n \geq 2$, $c_n^{1\text{-}2} = d_n^{1\text{-}2} = 2 \cdot 4^{n-2} + (n-2) \cdot 2^{n-2}$.

In the same manner, we can get that for the pattern 2-1,

$$\begin{cases} c_n^{2\text{-}1} = c_{n-1}^{2\text{-}1} + d_{n-1}^{2\text{-}1} + 2^{n-2} \cdot (2^{n-2}-1), \\ d_n^{2\text{-}1} = c_{n-1}^{2\text{-}1} + d_{n-1}^{2\text{-}1} + 2^{n-2} \cdot (2^{n-2}-1), \end{cases}$$

together with $c_3^{2\text{-}1} = 2$ and $d_3^{2\text{-}1} = 2$. This gives, that for all $n \geq 2$, $c_n^{2\text{-}1} = d_n^{2\text{-}1} = 2 \cdot 4^{n-2} - n \cdot 2^{n-2}$. □

Proposition 3 shows that asymptotically, the numbers of occurrences of the patterns, or the subsequences, 1-2 and 2-1 in C_n or D_n are equal.

4. PATTERNS WITHOUT INTERNAL DASHES

Recall the definitions in Section 2.

Theorem 4. *Let $\tau = \tau_1\tau_2\ldots\tau_\ell$ be an arbitrary generalized pattern without internal dashes that consists of 1s and 2s. Suppose $k = \lceil \log_2 \ell \rceil$, $a = m_\ell(\tau, D_k 1 C_k)$, and $b = m_\ell(\tau, D_k 2 C_k)$. Then for $n > k+1$, we have*

$$c_n^\tau = (a + b + c_{k+1}^\tau + d_{k+1}^\tau) \cdot 2^{n-k-2} - b,$$
$$d_n^\tau = (a + b + c_{k+1}^\tau + d_{k+1}^\tau) \cdot 2^{n-k-2} - a.$$

Proof. Suppose $n > k+1$. In this case, $C_n = C_{n-1} 1 D_{n-1} = W_1 \mathcal{K}_\ell(D_k 1 C_k) W_2$, for some words W_1 and W_2 such that $|W_1| = |W_2|$. Because of the definition of the kernel $\mathcal{K}_\ell(D_k 1 C_k)$, an occurrence of the pattern τ in C_n is in either C_{n-1}, or D_{n-1}, or $\mathcal{K}_k(D_k 1 C_k)$ (from the definitions $|C_{n-1} \cap \mathcal{K}_k(D_k 1 C_k)| = |D_{n-1} \cap \mathcal{K}_k(D_k 1 C_k)| = \ell - 1$ and thus these intersections cannot be an occurrence of τ). So,

$$c_n^\tau = c_{n-1}^\tau + d_{n-1}^\tau + a,$$

whereas in the same way, we can obtain that

$$d_n^\tau = c_{n-1}^\tau + d_{n-1}^\tau + b.$$

By solving these recurrence relations, we get the desirable result. □

In particular, Theorem 4 is valid for $\ell = 1$, in which case the number of occurrences of τ in C_n (or D_n) is the number of letters in C_n (or D_n). Indeed, in this case, $k = 0$, $a = b = c_1^1 = d_1^1 = 1$, hence $c_n^1 = d_n^1 = 2^n - 1 = |C_n| = |D_n|$. Also, as a corollary to Theorem 4 we have, that if $a = b = c_{k+1}^\tau = d_{k+1}^\tau = 0$ for some pattern τ, then this pattern never appears in sigma-sequence.

All of the following examples are corollaries to Theorem 4.

Example 5. Suppose $\tau = 12$. We have that $k = 1$, $a = m_2(12, D_1 1C_1) = 0$ and $b = m_2(12, D_1 2C_1) = 0$. Besides, $c_2^{12} = 1$ and $d_2^{12} = 1$. Thus using Theorem 4, for all $n > 2$, $c_n^{12} = 2^{n-2}$. So, the number of rises in C_n is equal to 2^{n-2}, for $n \geq 2$.

If $\tau = 21$, again $k = 1$, but now $a = m_2(21, D_1 1C_1) = 1$ and $b = m_2(21, D_1 2C_1) = 1$. Besides, $c_3^{21} = 1$ and $d_3^{21} = 1$. From Theorem 4, for all $n > 3$, $c_n^{21} = 2^{n-2} - 1$, which shows that the number of descents in C_n is one less than the number of rises.

Since in both cases $a = b$, using the recurrences in Theorem 4, we have that $c_n^{12} = d_n^{12} = 2^{n-2}$, whereas $c_n^{21} = d_n^{21} = 2^{n-2} - 1$.

Example 6. Suppose $\tau = 112$. We have that $k = 2$, $a = m_3(112, D_2 1C_2) = 0$, and $b = m_3(112, D_2 2C_2) = 0$. Besides, $c_3^{112} = 2$ and $d_3^{112} = 1$. Now, from Theorem 4, we have that for all $n > 3$, $c_n^{112} = d_n^{112} = 3 \cdot 2^{n-4}$.

Example 7. Suppose $\tau = 221$. We have that $k = 2$, $a = m_3(221, D_2 1C_2) = 1$, and $b = m_3(221, D_2 2C_2) = 1$. Besides, $c_3^{221} = 0$ and $d_3^{221} = 1$. Now, from Theorem 4, we have that for all $n > 3$, $c_n^{221} = d_n^{221} = 3 \cdot 2^{n-4} - 1$.

Example 8. If $\tau = 2212221$ then $k = 3$, $a = m_7(221, D_3 1C_3) = 0$, $b = m_7(221, D_3 2C_3) = 1$, $c_4^{2212221} = 0$, and $d_4^{2212221} = 0$. Thus for $n \geq 4$, $c_n^{2212221} = 2^{n-4} - 1$.

5. PATTERNS OF THE FORM τ_1-τ_2

Theorem 9. *Let $p = \tau_1$-τ_2 be a generalized pattern such that $|\tau_1| = k_1$ and $|\tau_2| = k_2$. Suppose $k = \lceil \log_2(k_1 + k_2 - 1) \rceil$. The following denote the number of occurrences of the subwords τ_1 and τ_2 in the certain kernels: $a_{\tau_1} = m_{k_1}(\tau_1, D_k 1C_k)$, $a_{\tau_2} = m_{k_2}(\tau_2, D_k 1C_k)$, $b_{\tau_1} = m_{k_1}(\tau_1, D_k 2C_k)$, and $b_{\tau_2} = m_{k_2}(\tau_2, D_k 2C_k)$. Also, let r_1^a (resp. r_2^a, r_1^b, r_2^b) denote the number of occurrences of overlapping subwords τ_1 and τ_2 in the word $D_k 1C_k$ (resp. $D_k 1C_k$, $D_k 2C_k$, $D_k 2C_k$), where $\tau_1 \in \mathcal{K}_{k_1}(D_k 1C_k)$ and $\tau_2 \in C_k$ (resp. $\tau_1 \in D_k$ and $\tau_2 \in \mathcal{K}_{k_2}(D_k 1C_k)$, $\tau_1 \in \mathcal{K}_{k_1}(D_k 2C_k)$ and $\tau_2 \in C_k$, $\tau_1 \in D_k$ and $\tau_2 \in \mathcal{K}_{k_2}(D_k 2C_k)$). Besides, we assume that we know $c_n^{\tau_i}$ and $d_n^{\tau_i}$ for $n > n_i$, $i = 1, 2$. Then for $n > \max(k+1, n_1+1, n_2+1)$, c_n^{τ} and d_n^{τ} are given by the following recurrence:*

$$\begin{pmatrix} c_n^{\tau} \\ d_n^{\tau} \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} c_{n-1}^{\tau} \\ d_{n-1}^{\tau} \end{pmatrix} + \begin{pmatrix} \alpha_n \\ \beta_n \end{pmatrix},$$

where

$$\alpha_n = (c_{n-1}^{\tau_1} + a_{\tau_1} - r_1^a) d_{n-1}^{\tau_2} + (a_{\tau_2} - r_2^a) c_{n-1}^{\tau_1}$$

and

$$\beta_n = (c_{n-1}^{\tau_1} + b_{\tau_1} - r_1^b) d_{n-1}^{\tau_2} + (b_{\tau_2} - r_2^b) c_{n-1}^{\tau_1}.$$

Proof. Suppose $n > \max(k+1, n_1+1, n_2+1)$. Let us find a recurrence for the number c_n^{τ} (one can use the same considerations for d_n^{τ}).

An occurrence of the pattern τ in $C_n = C_{n-1} 1 D_{n-1}$ is either inside C_{n-1}, or inside D_{n-1}, or begins in C_{n-1} or the letter 1 between C_{n-1} and D_{n-1} and ends in D_{n-1} or the letter 1. The first two cases obviously give c_{n-1} and d_{n-1} occurrences of τ. To count the contribution of the last to cases, we work with words instead of patterns. We do it to take in account the situations when τ_1 or τ_2 consists of copies of only one letter. In this case, we cannot count occurrence of these patterns separately, and then use this information, since, for instance, occurrences of the pattern $\tau_1 = 111$ are subwords 111 and 222 (the last one of these subwords we do not need), whereas occurrences of the pattern $\tau_1 = 222$ are not defined at all (222 is not a pattern).

If an occurrence of τ_1-τ_2 does not entirely belong to C_{n-1} or D_{n-1} then we only have one of the following possibilities:

(a) the subword τ_1 entirely belongs to C_{n-1} and the subword τ_2 entirely belongs to D_{n-1};
(b) the subword τ_1 belongs entirely to C_{n-1} and the subword τ_2 belongs to the kernel $\mathcal{K}_{k_2}(D_k 1C_k)$, where $k = \lceil \log_2(k_1+k_2-1) \rceil$ is the least number that allow to control, in C_n ($n > k$), overlapping occurrences of subwords τ_1 and τ_2 where τ_1 is entirely from C_{n-1} and $\tau_2 \in \mathcal{K}_{k_2}(D_k 1C_k)$;
(c) the subword τ_2 belongs entirely to D_{n-1} and the subword τ_1 belongs to the kernel $\mathcal{K}_{k_1}(D_k 1C_k)$.

In (a) we obviously have $c^{\tau_1}_{n-1} \cdot d^{\tau_2}_{n-1}$ possibilities.

In (b) we have $c^{\tau_1}_{n-1} \cdot a_{\tau_2} - c^{\tau_1}_{n-1} \cdot r^a_2$ possibilities, since we need to subtract those occurrences of τ_1 and τ_2 that overlap.

Similarly to (b), in (c) we have $d^{\tau_2}_{n-1} \cdot a_{\tau_1} - d^{\tau_2}_{n-1} \cdot r^a_1$ possibilities, which completes the proof. □

Remark 10. For using Theorem 9, one needs to know c^τ_n and d^τ_n for patterns τ without internal dashes. These numbers could be obtained by using Theorem 4.

The following corollary to Theorem 9 is straitforward to prove, using the fact that for non-overlapping patterns τ_1 and τ_2, $r^a_1 = r^a_2 = r^b_1 = r^b_2 = 0$.

Corollary 11. *We make the same assumptions as those in Theorem 9. Suppose additionally that the words* τ_1 *and* τ_2 *are not overlapping in the following sense: no one suffix of* τ_1 *is a prefix of* τ_2*. Then for* $n > \max(k+1, n_1+1, n_2+1)$*,* c^τ_n *and* d^τ_n *are given by the same recurrence as that in Theorem 9 with*

$$\alpha_n = (c^{\tau_1}_{n-1} + a_{\tau_1})d^{\tau_2}_{n-1} + a_{\tau_2}c^{\tau_1}_{n-1}$$

and

$$\beta_n = (c^{\tau_1}_{n-1} + b_{\tau_1})d^{\tau_2}_{n-1} + b_{\tau_2}c^{\tau_1}_{n-1}.$$

Remark 12. Corollary 11 is valid under weaker assumptions, namely we only need the property of non-overlapping of the patterns τ_1 and τ_2 when one of them is in its kernel and the other one is not in its kernel. Example 15 deals with the pattern τ that has overlapping blocks τ_1 and τ_2, but Corollary 11 can be applied. However, from practical point of view, checking the fact if two subwords are non-overlapping is more easy than considering the kernels and checking the non-overlapping of the subwords there.

Example 13. Suppose $\tau = $ 12-21. We have that $|\tau_1| = |\tau_2| = 2$. Now, in the statement of Theorem 9 we have that $k = 2$, $a_{\tau_1} = 0$, $a_{\tau_2} = 1$, $b_{\tau_1} = 0$ and $b_{\tau_2} = 1$. Also, since there are no overlapping occurrences of the subwords 12 and 21 in $\mathcal{K}_3(1221112)$ and $\mathcal{K}_3(1222112)$, we have $r^a_1 = 0$, $r^a_2 = 0$, $r^b_1 = 0$ and $r^b_2 = 0$. Besides, from example 5, $c^{12}_n = d^{12}_n = 2^{n-2}$ and $c^{21}_n = d^{21}_n = 2^{n-2} - 1$. Thus, $\alpha_n = \beta_n = 4^{n-3}$. Using the fact that $c^{12\text{-}21}_3 = 0$ and $d^{12\text{-}21}_3 = 1$, this allows us to get an explicit formula for $c^{12\text{-}21}_n$ and $d^{12\text{-}21}_n$ for $n > 3$:

$$c^{12\text{-}21}_n = d^{12\text{-}21}_n = \frac{1}{2}4^{n-2} - 3 \cdot 2^{n-4}.$$

In particular $c^{12\text{-}21}_4 = 5$.

Example 14. Suppose $\tau = $ 1-221. We have that $|\tau_1| = 1$ and $|\tau_2| = 3$. Moreover, the words τ_1 and τ_2 are not overlapping, hence we can use Corollary 11. We have that $k = 2$, $a_{\tau_1} = 1$, $a_{\tau_2} = 1$, $b_{\tau_1} = 0$ and $b_{\tau_2} = 1$. From example 7, $d^{221}_n = 3 \cdot 2^{n-4} - 1$. Also, the number of occurrences of the letter 1 (the subword $\tau_1 = 1$) is given by Lemma 1: $c^1_n = 2^{n-1}$. So, $\alpha_n = 6 \cdot 4^{n-4} + 3 \cdot 2^{n-5} - 1$ and $\beta_n = 6 \cdot 4^{n-4}$. One can get now an explicit formula for $c^{1\text{-}221}_n$ and $d^{1\text{-}221}_n$ for $n > 4$:

$$c^{1\text{-}221}_n = \tfrac{1}{2}4^{n-2} + 27 \cdot 2^{n-5} - n - 7,$$
$$d^{1\text{-}221}_n = \tfrac{1}{2}4^{n-2} + 21 \cdot 2^{n-5} - 8.$$

In particular, $c^{1\text{-}221}_5 = 47$.

Example 15. Suppose $\tau = $ 112-21. We have that $|\tau_1| = k_1 = 3$ and $|\tau_2| = k_2 = 2$. The other parameters in Theorem 9 are $k = 3$, $a_{\tau_1} = 0$, $a_{\tau_2} = 1$, $b_{\tau_1} = 0$, $b_{\tau_2} = 1$, $r^a_1 = r^a_2 = r^b_1 = r^b_2 = 0$. From Example 6, for $n \geq 4$, $c^{112}_n = 3 \cdot 2^{n-4}$, and from Example 5, $d^{21}_n = 2^{n-2} - 1$. Thus, in Theorem 9, $\alpha_n = \beta_n = c^{112}_{n-1}(d^{21}_{n-1} + 1) = 3 \cdot 4^{n-4}$. Now, we solve the recurrence relation from the theorem to get, that for $n > 3$

$$c^{112\text{-}21}_n = d^{112\text{-}21}_n = \frac{3}{2} \cdot 4^{n-3} - 2^{n-4}.$$

6. COUNTING OCCURRENCES OF $\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k$

In this section we study the number of occurrences of a pattern $\tau = \tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k$, where τ_i are patterns without internal dashes. We say that τ consists of k blocks. We assume that for $i = 1, 2, \ldots, k-1$, the pattern τ_i does not overlap with the patterns τ_{i-1} and τ_{i+1}. In this case we give a recurrence relation for the number of occurrences of τ, provided we know the number of occurrences of certain patterns consisting of less than, or equal to, $k-1$ blocks, as well as $2k$ certain numbers which can be calculated by considering the words $D_\ell 1 C_\ell$ and $C_\ell 2 D_\ell$, where ℓ is the maximum number such that $\ell \leq \max_i \lceil \log_2 |\tau_i| \rceil$. The cases of $k = 1$ and $k = 2$ are studied in the previous sections; they give the bases for our calculations. However, the case of overlapping patterns τ_i is not solved, and it remains as a challenging problem, since an answer to this problem gives the way to count occurrences of an arbitrary generalized pattern, or an arbitrary subsequence, in σ-words.

Theorem 16. *Let $\tau = \tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k$ be a generalized pattern such that $|\tau_i| = k_i$ for $i = 1, 2, \ldots, k$. We assume that for $i = 1, 2, \ldots, k-1$, the subword τ_i does not overlap with the subwords τ_{i-1} and τ_{i+1} in the following sense: no one suffix of τ_{i-1} is a prefix of τ_i and no one suffix of τ_i is a prefix of τ_{i+1}. Suppose $\ell_i = \lceil \log_2 k_i \rceil$, $\ell = \max_i \ell_i$, and for the subwords τ_i we have $a_i = m_{k_i}(\tau_i, D_{\ell_i} 1 C_{\ell_i})$ and $b_i = m_{k_i}(\tau_i, D_{\ell_i} 2 C_{\ell_i})$, for $i = 1, 2, \ldots, k$. We assume that we know $c_{n-1}^{\tau_1\text{-}\cdots\text{-}\tau_i}$ and $d_{n-1}^{\tau_{i+1}\text{-}\cdots\text{-}\tau_k}$ for each $1 \leq i \leq k-1$ and for all $n > n^\star$. Then for all $n > \max(\ell+1, n^\star+1)$, c_n^τ and d_n^τ are given by the following recurrence:*

$$\begin{pmatrix} c_n^\tau \\ d_n^\tau \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} c_{n-1}^\tau \\ d_{n-1}^\tau \end{pmatrix} + \sum_{i=1}^{k-1} \begin{pmatrix} c_{n-1}^{\tau_1\text{-}\cdots\text{-}\tau_i} \cdot d_{n-1}^{\tau_{i+1}\text{-}\cdots\text{-}\tau_k} \\ c_{n-1}^{\tau_1\text{-}\cdots\text{-}\tau_i} \cdot d_{n-1}^{\tau_{i+1}\text{-}\cdots\text{-}\tau_k} \end{pmatrix} + \sum_{i=1}^{k} \begin{pmatrix} a_i \cdot c_{n-1}^{\tau_1\text{-}\cdots\text{-}\tau_{i-1}} \cdot d_{n-1}^{\tau_{i+1}\text{-}\cdots\text{-}\tau_k} \\ b_i \cdot c_{n-1}^{\tau_1\text{-}\cdots\text{-}\tau_{i-1}} \cdot d_{n-1}^{\tau_{i+1}\text{-}\cdots\text{-}\tau_k} \end{pmatrix}.$$

Proof. We consider only c_n^τ, since the same arguments can be applied to d_n^τ. We use the considerations similar to those in Theorem 9.

An occurrence of the pattern τ in $C_n = C_{n-1} 1 D_{n-1}$ can be entirely in C_n or D_n. The first term counts such occurrences. Otherwise, we have two possibilities: either the letter 1 between the words C_{n-1} and D_{n-1} does not belong to an occurrence of τ, or it does do it, in which case there exist i (exactly one) such that the subword τ_i occurs in its kernel. The first sum in the statement is obviously responsible for the first of this cases, whereas the second sum is responsible for the second case (in the last case we use the fact that subwords τ_i are not overlapping). □

As a corollary to Theorem 16, we have Corollary 11.

The following example is another corollary to Theorem 16.

Example 17. Suppose $\tau = 2 - 1 - 221$, that is, $\tau_1 = 2$, $\tau_2 = 1$ and $\tau_3 = 221$. So, parameters in Theorem 16 are the following: $k_1 = k_2 = 1$, $k_3 = 3$, $\ell_1 = \ell_2 = 1$, $\ell_3 = 2$, $\ell = 2$. From $D_1 1 C_1 = 211$ we obtain $a_1 = 0$, $a_2 = 1$. From $D_2 1 C_2 = 1221112$ we obtain $a_3 = 1$. From $D_1 2 C_1 = 221$ we get $b_1 = 1$, $b_2 = 0$. From $D_2 2 C_2 = 1222112$ we get $b_3 = 1$. Besides, from Proposition 3, Examples 7 and 14, we have

$$c_n^{\tau_1\text{-}\tau_2} = c_n^{2\text{-}1} = 2 \cdot 4^{n-2} - n \cdot 2^{n-2}, \text{for } n > 1;$$

$$d_n^{\tau_3} = d_n^{221} = 3 \cdot 2^{n-4} - 1, \text{for } n > 3;$$

$$d_n^{\tau_2\text{-}\tau_3} = d_n^{1\text{-}221} = \tfrac{1}{2} \cdot 4^{n-2} + 21 \cdot 2^{n-5} - 8, \text{for } n > 4.$$

Also, the number of occurrences of the subword $\tau_1 = 2$ in C_n is given by Proposition 2: $c_n^{\tau_1} = c_n^2 = 2^{n-1} - 1$. So, the number of occurrences of the pattern τ in C_n and D_n, for $n > 5$, satisfies the following recurrence relation:

$$\begin{pmatrix} c_n^\tau \\ d_n^\tau \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} c_{n-1}^\tau \\ d_{n-1}^\tau \end{pmatrix} + \begin{pmatrix} \frac{5}{1024} 8^n + \frac{25-3n}{256} 4^n - \frac{171}{64} 2^n + 9 \\ \frac{5}{1024} 8^n + \frac{21-3n}{256} 4^n - 2^{n+1} \end{pmatrix},$$

with initial conditions $c_5^\tau = 70$ and $d_5^\tau = 74$.

7. PATTERNS OF THE FORM $[\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k]$, $[\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k)$ AND $(\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k]$

We recall that according to Babson and Steingrímsson notation for generalized patterns, if we use "[" in a pattern, for example if we write $p = [1\text{-}2)$, we indicate that in an occurrence of p, the letter corresponding to the 1 must be the first letter of a word under consideration, whereas if we write, say, $p = (1\text{-}2]$, then the letter corresponding to 2 must be the last (rightmost) letter of the word.

In the theorems of this section, we assume that we can find the numbers $c_n^{\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k}$ and $d_n^{\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k}$ for any patterns τ_i, $i = 1, 2, \dots, k$, without internal dashes. For certain special cases, these numbers can be obtained using the theorems of Sections 5 and 6.

Theorem 18. *Suppose τ_1 and τ_2 are two patterns without internal dashes such that $|\tau_1| = k_1$ and $|\tau_2| = k_2$. Also, suppose $\ell_1 = \log_2(k_1+1)$, $\ell_2 = \log_2(k_2+1)$ and $\ell = \log_2(k_1+k_2+1)$. Let $a(\tau_1,\tau_2)$ be the number of overlapping subwords τ_1 and τ_2 in C_ℓ such that τ_1 occurs as k_1 leftmost letters of C_ℓ; $b(\tau_1,\tau_2)$ is the number of overlapping subwords τ_1 and τ_2 in C_ℓ such that τ_2 occurs as k_2 rightmost letters of C_ℓ. We assume that we know $c_n^{\tau_i}$ and $d_n^{\tau_i}$ for $i = 1, 2$ and for all $n > n^\star$.*

i. *For $n \geq \max(\ell_1, n^\star)$,*

$$c_n^{[\tau_1\text{-}\tau_2)} = \begin{cases} 0, & \text{if } C_{\ell_1} \text{ does not begin with } \tau_1, \\ c_n^{\tau_2} - a(\tau_1,\tau_2), & \text{otherwise.} \end{cases}$$

ii. *For $n \geq \max(\ell_2, n^\star)$,*

$$c_n^{(\tau_1\text{-}\tau_2]} = \begin{cases} 0, & \text{if } C_{\ell_2} \text{ does not end with } \tau_2, \\ c_n^{\tau_1} - b(\tau_1,\tau_2), & \text{otherwise.} \end{cases}$$

iii. *For $n \geq \ell$,*

$$c_n^{[\tau_1\text{-}\tau_2]} = \begin{cases} 0, & \text{if } C_\ell \text{ does not begin with } \tau_1 \text{ or end with } \tau_2, \\ 1, & \text{otherwise.} \end{cases}$$

iv. *For $n \geq \max(\ell_1, n^\star)$,*

$$d_n^{[\tau_1\text{-}\tau_2)} = \begin{cases} 0, & \text{if } D_{\ell_1} \text{ does not begin with } \tau_1, \\ d_n^{\tau_2} - a(\tau_1,\tau_2), & \text{otherwise.} \end{cases}$$

v. *For $n \geq \max(\ell_2, n^\star)$,*

$$d_n^{(\tau_1\text{-}\tau_2]} = \begin{cases} 0, & \text{if } D_{\ell_2} \text{ does not end with } \tau_2, \\ d_n^{\tau_1} - b(\tau_1,\tau_2), & \text{otherwise.} \end{cases}$$

vi. *For $n \geq \ell$,*

$$d_n^{[\tau_1\text{-}\tau_2]} = \begin{cases} 0, & \text{if } D_\ell \text{ does not begin with } \tau_1 \text{ or end with } \tau_2, \\ 1, & \text{otherwise.} \end{cases}$$

Proof. We prove case i, all the other cases are then easy to see.

Clearly, if C_{ℓ_1} does not begin with τ_1 then C_n does not begin with τ_1 for all $n \geq \ell_1$, which means that $c_n^{[\tau_1\text{-}\tau_2)} = 0$ in this case. Otherwise, to count occurrences of the pattern $[\tau_1\text{-}\tau_2)$ is the same as to find the number of occurrences of the pattern τ_2 in C_n and then subtract the number of such occurrences of τ_2 that begin from the i-th letter of C_n, where $1 \leq i \leq k_1$. □

The following two examples are corollaries to Theorem 18.

Example 19. Suppose we have the patterns $\sigma_1 = [1122 - 21211)$ and $\sigma_2 = (21221 - 12]$. From Theorem 18, $c_n^{\sigma_1} = d_n^{\sigma_1} = 0$, since C_3 does not begin with 1122 ($\ell_1 = 3$). Also, $c_n^{\sigma_2} = d_n^{\sigma_2} = 0$, since C_3 does not end with 12 ($\ell_2 = 3$).

Example 20. Suppose $\tau = [112\text{-}21)$. We have that $k_1 = 3$, $\ell_1 = 2$ and C_2 begins with the subword 112. Besides, $a(112, 21) = 1$ and, from Example 5, $c_n^{21} = d_n^{21} = 2^{n-2} - 1$. Theorem 18 now yields, that for $n > 3$, we have $c_n^{[112\text{-}21)} = c_n^{\tau_2} - a(\tau_1,\tau_2) = 2^{n-2} - 2$.

The following theorem is straightforward to prove under the assumption that certain subwords do not overlap.

Theorem 21. *Let* $\{\tau_1, \tau_2, \ldots, \tau_k\}$ *be a set of generalized patterns without internal dashes. Suppose* $|\tau_1| = s_1$*,* $|\tau_k| = s_k$*,* $\ell_1 = \log_2(s_1+1)$ *and* $\ell_k = \log_2(s_k+1)$*. Also,* $\ell = \max(\ell_1, \ell_k)$*.*

i. *With the assumption that the subword* τ_1 *does not overlap with the subword* τ_2*, that is, no one suffix of* τ_1 *is a prefix of* τ_2*, we have*

(a)

$$c_n^{[\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k)} = \begin{cases} 0, & \textit{if } C_{\ell_1} \textit{ does not begin with } \tau_1, \\ c_n^{\tau_2\text{-}\tau_3\text{-}\cdots\text{-}\tau_k}, & \textit{otherwise.} \end{cases}$$

(b)

$$d_n^{[\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k)} = \begin{cases} 0, & \textit{if } D_{\ell_1} \textit{ does not begin with } \tau_1, \\ d_n^{\tau_2\text{-}\tau_3\text{-}\cdots\text{-}\tau_k}, & \textit{otherwise.} \end{cases}$$

ii. *With assumption that the subword* τ_{k-1} *does not overlap with the subword* τ_k*, that is, no one suffix of* τ_{k-1} *is a prefix of* τ_k*, we have*

(a)

$$c_n^{(\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k]} = \begin{cases} 0, & \textit{if } C_{\ell_k} \textit{ does not end with } \tau_k, \\ c_n^{\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_{k-1}}, & \textit{otherwise.} \end{cases}$$

(b)

$$d_n^{(\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k]} = \begin{cases} 0, & \textit{if } D_{\ell_k} \textit{ does not end with } \tau_k, \\ d_n^{\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_{k-1}}, & \textit{otherwise.} \end{cases}$$

iii. *With the assumption that the subword* τ_1 *does not overlap with the subword* τ_2*, and the subword* τ_{k-1} *does not overlap with the subword* τ_k*, we have*

(a)

$$c_n^{[\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k]} = \begin{cases} 0, & \textit{if } C_{\ell} \textit{ does not begin with } \tau_1 \textit{ or does not end with } \tau_k, \\ c_n^{\tau_2\text{-}\tau_3\text{-}\cdots\text{-}\tau_{k-1}}, & \textit{otherwise.} \end{cases}$$

(b)

$$d_n^{[\tau_1\text{-}\tau_2\text{-}\cdots\text{-}\tau_k]} = \begin{cases} 0, & \textit{if } D_{\ell} \textit{ does not begin with } \tau_1 \textit{ or does not end with } \tau_k, \\ d_n^{\tau_2\text{-}\tau_3\text{-}\cdots\text{-}\tau_{k-1}}, & \textit{otherwise.} \end{cases}$$

The following example is a corollary to Theorem 21.

Example 22. Suppose $\tau = [112\text{-}1\text{-}221\text{-}22]$. The parameters of Theorem 21 are $k_1 = 3$, $k_2 = 2$, $\ell_1 = 2$, $\ell_2 = 2$, $\ell = 2$. C_3 begins with the subword 112 and ends with the subword 22. Thus by Theorem 21 and Example 14, $c_n^{[112\text{-}1\text{-}221\text{-}22]} = c_n^{1\text{-}221} = \frac{1}{2}4^{n-2} + 27\cdot 2^{n-5} - n - 7$.

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